The "Second Law" of Probability: Entropy Growth in the Central Limit Theorem.

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The second law of thermodynamics

Joule and Carnot studied ways to improve the efficiency of steam engines.

Is it possible for a thermodynamic system to move from state A to state B without any net energy being put into the system from outside?

A single experimental quantity, dubbed **entropy**, made it possible to decide the direction of thermodynamic changes.

The second law of thermodynamics

The entropy of a closed system increases with time.

The second law applies to all changes: not just thermodynamic.

Entropy measures the extent to which energy is dispersed: so the second law states that energy tends to disperse.

The second law of thermodynamics



Closed systems become progressively more featureless.

We expect that a closed system will approach an equilibrium with maximum entropy.

Information theory

Shannon showed that a noisy channel can communicate information with almost perfect accuracy, up to a fixed rate: the capacity of the channel.

The (Shannon) entropy of a probability distribution: if the possible states have probabilities $p_1, p_2, ..., p_n$ then the entropy is $p_i \log_2 p_i$

Entropy measures the number of (YES/NO) questions that you expect to have to ask in order to find out which state has occurred.

Information theory

You can distinguish 2^k states with *k* (YES/NO) questions.

If the states are equally likely, then this is the best you can do.



It costs k questions to identify a state from among 2^k equally likely states.

Information theory

It costs k questions to identify a state from among 2^k equally likely states.

It costs $\log_2 n$ questions to identify a state from among *n* equally likely states: to identify a state with probability 1/n.

Probability	Questions
1/ <i>n</i>	$\log_2 n$
p	$\log_2(1/p)$

The entropy

State	Probability	Questions	Uncertainty
S_1	p_1	$\log_2(1/p_1)$	$p_1 \log_2 (1/p_1)$
<i>S</i> ₂	p_2	$\log_2(1/p_2)$	$p_2 \log_2 (1/p_2)$
S_3	p_3	$\log_2(1/p_3)$	$p_3 \log_2(1/p_3)$

Entropy = $p_1 \log_2(1/p_1) + p_2 \log_2(1/p_2) + p_3 \log_2(1/p_3) + \dots$

Continuous random variables

For a random variable X with density f the entropy is

$$\operatorname{Ent}(X) = ! \quad f \log f$$

The entropy behaves nicely under several natural processes: for example, the evolution governed by the heat equation.

If the density f measures the distribution of heat in an infinite metal bar, then f evolves according to the heat equation:

$$\frac{!f}{!t} = f"$$

The entropy increases:

$$\frac{!}{!t} \left(f \log f \right) = \frac{f'^2}{f} \# 0$$

Fisher information

The central limit theorem

If X_i are independent copies of a random variable with mean 0 and finite variance, then the normalized sums

$$\frac{1}{\sqrt{n}} \prod_{i=1}^{n} X_{i}$$

converge to a Gaussian (normal) with the same variance.

Most proofs give little intuition as to why.

The central limit theorem

Among random variables with a given variance, the Gaussian has largest entropy.

Theorem (Shannon-Stam) If *X* and *Y* are independent and identically distributed, then the normalized sum

 $\frac{X+Y}{\sqrt{2}}$

has entropy at least that of *X* and *Y*.

Idea

The central limit theorem is analogous to the second law of thermodynamics: the normalized sums

$$S_n = \frac{1}{\sqrt{n}} \prod_{i=1}^n X_i$$

have increasing entropy which drives them to an "equilibrium" which has maximum entropy.

Problem: (folklore or Lieb (1978)).

Is it true that $Ent(S_n)$ increases with *n*?

Shannon-Stam shows that it increases as n goes from 1 to 2 (hence 2 to 4 and so on). Carlen and Soffer found uniform estimates for entropy jump from 1 to 2.

It wasn't known that entropy increases from 2 to 3.

The difficulty is that you can't express the sum of 3 independent random variables in terms of the sum of 2: you can't add 3/2 independent copies of *X*.

The Fourier transform?

The simplest proof (conceptually) of the central limit theorem uses the FT. If X has density f whose FT is ! then the FT of the density of $! X_i$ is $!^n$.

The problem is that the entropy cannot easily be expressed in terms of the FT. So we must stay in real space instead of Fourier space.

Example:

Suppose *X* is uniformly distributed on the interval between 0 and 1. Its density is:

When we add two copies the density is:



For 9 copies the density is a spline defined by 9 different polynomials on different parts of the range.



The central polynomial (for example) is:

3 (857291 - 5027400 x + 12800340 x² - 18438840 x³ + 16391970 x⁴ - 9185400 x⁵ + 3163860 x⁶ - 612360 x⁷ + 51030 x⁸)/4480

and its logarithm is?

The second law of probability

A new variational approach to entropy gives quantitative measures of entropy growth and proves the "second law".

Theorem (Artstein, Ball, Barthe, Naor) If X_i are independent copies of a random variable with finite variance, then the normalized sums

$$\frac{1}{\sqrt{n}} \prod_{i=1}^{n} X_{i}$$

have increasing entropy.

Starting point: used by many authors. Instead of considering entropy directly, we study the Fisher information:

$$J(X) = \int \frac{f'^2}{f}$$

Among random variables with variance 1, the Gaussian has the **smallest** Fisher information, namely 1.

The Fisher information should decrease as a process evolves.

The connection (we want) between entropy and Fisher information is provided by the Ornstein-Uhlenbeck process (de Bruijn, Bakry and Emery, Barron).

Recall that if the density of $X^{(t)}$ evolves according to the heat equation then

$$\frac{!}{!t}$$
Ent $(X^{(t)}) = J(X^{(t)})$

The heat equation can be solved by running a Brownian motion from the initial distribution. The Ornstein-Uhlenbeck process is like Brownian motion but run in a potential which keeps the variance constant.

The Ornstein-Uhlenbeck process

A discrete analogue:



You have *n* sites, each of which can be ON or OFF. At each time, pick a site (uniformly) at random and switch it.

 $X^{(t)} =$ (number on)-(number off).

The Ornstein-Uhlenbeck process



The Ornstein-Uhlenbeck evolution

The density evolves according to the modified diffusion equation:

$$\frac{!f}{!t} = f"+(xf)'$$

From this:

$$\frac{!}{!t}$$
Ent $(X^{(t)}) = J(X^{(t)})$ 1

As $t \,!$ the evolutes approach the Gaussian of the same variance.

The entropy gap can be found by integrating the information gap along the evolution.

Ent(G) Ent(
$$X^{(0)}$$
) = $\# (J(X^{(t)}) \ 1)$

In order to prove entropy increase, it suffices to prove that the information

$$J(n) = J \frac{!}{\# \sqrt{n}} \prod_{i=1}^{n} X_{i}$$

decreases with *n*.

It was known (Blachman-Stam) that $J(2) \downarrow J(1)$.

Main new tool: a variational description of the information of a marginal density.

If w is a density on ⁿ and e is a unit vector, then the marginal in direction e has density



Main new tool:

The density *h* is a marginal of *w* and

$$J(h) = \frac{h'(t)^2}{h(t)} = \frac{h'(t)^2}{h(t)}! h''(t) = h(! \log h)''$$

The integrand is non-negative if *h* has concave logarithm.

Densities with concave logarithm have been widely studied in high-dimensional geometry, because they naturally generalize convex solids.

The Brunn-Minkowski inequality



Let A(x) be the cross-sectional area of a convex body at position x.

Then log A is concave.

The function *A* is a marginal of the body.

The Brunn-Minkowski inequality



We can replace the body by a function with concave logarithm. If *w* has concave logarithm, then so does each of its marginals.

If the density h is a marginal of w, the inequality tells us something about $(! \log h)$ " in terms of Hess $(! \log w)$

The Brunn-Minkowski inequality

If the density h is a marginal of w, the inequality tells us something about $(! \log h)$ " in terms of Hess $(! \log w)$

We rewrite a proof of the Brunn-Minkowski inequality so as to provide an explicit relationship between the two. The expression involving the Hessian is a quadratic form whose minimum is the information of h.

This gives rise to the variational principle.

The variational principle

Theorem If *w* is a density and *e* a unit vector then the information of the marginal in the direction *e* is

$$J(h) = \int \frac{h'(t)^2}{h(t)} dt = \min \int_{n}^{\infty} \frac{\operatorname{div}(pw)^2}{w}$$

where the minimum is taken over vector fields p satisfying !p,e #1.

$$J(h) = \prod_{h=1}^{n} \frac{h'(t)^2}{h(t)} dt = \min_{h=1}^{n} \frac{\operatorname{div}(pw)^2}{w}$$

Technically we have gained because h(t) is an integral: not good in the denominator.

The real point is that we get to choose p. Instead of choosing the optimal p which yields the intractable formula for information, we choose a non-optimal p with which we can work.

Proof of the variational principle.

$$h(t) = \lim_{t \in +\langle e \rangle^!} w$$

SO

$$h'(t) = \#_{e+\langle e \rangle^!} \quad {}_e W$$

If p satisfies !p, e #1 at each point, then we can realise the derivative as

$$h'(t) = \underset{te+\langle e \rangle^!}{\operatorname{div}(pw)}$$

since the part of the divergence perpendicular to *e* integrates to 0 by the Gauss-Green (divergence) theorem.

Hence

$$Q_{Q(t)}^{h'(t)^{2}} dt = Q_{Q(t)}^{(t)} \frac{div(pw)}{dt} dt = Q_{Q(t)}^{(t)} \frac{div(pw)}{dt} dt = Q_{Q(t)}^{(t)} \frac{div(pw)}{w}^{2}$$

There is equality if

$$\operatorname{div}(pw) = \frac{h'(t)}{h(t)}w$$

This divergence equation has many solutions: for example we might try the electrostatic field solution. But this does not decay fast enough at infinity to make the divergence theorem valid.

$$\operatorname{div}(pw) = \frac{h'(t)}{h(t)}w$$

The right solution for p is a flow in the direction of e which transports between the probability measures induced by w on hyperplanes perpendicular to e.

For example, if w is 1 on a triangle and 0 elsewhere, the flow is as shown. (The flow is irrelevant where w = 0.)



The second law of probability

Theorem If X_i are independent copies of a random variable with variance, then the normalized sums

$$\frac{1}{\sqrt{n}} \prod_{i=1}^{n} X_{i}$$

have increasing entropy.